

Safe Composition of CFSM Systems via Partial Gateways

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University of Catania

ICE'26 - June 12, 2026, Urbino

OVERVIEW

- ▶ Introduction: the need for system composability
- ▶ The “participants-as-interfaces” (PaI) approach to **binary** system composition
[BARBANERA-HENNICKER-DE’LIGUORO]
- ▶ From binary PaI to:
 - multicomposition [BARBANERA-HENNICKER];
 - orchestrated multicomposition [BARBANERA-HENNICKER];
 - composition via partial gateways **ICE’26**;
- ▶ Exploiting the PaI approach in the asynchronous framework of Communicating **Finite State Machines**: Safety of the approach.
- ▶ The essential of PaI: (binary) partial-fusion composition **ICE’26**.

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The need of systems composability

- ▶ Concurrent/Distributed systems are
not STAND-ALONE ENTITIES
- ▶ (especially nowadays) they are parts of
JIGSAWS NEVER COMPLETELY TERMINATED

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Good composition methods

They should be

They should be

▶ CONSERVATIVE

Altering as less as possible the single systems



They should be

▶ CONSERVATIVE

▶ FLEXIBLE

- i.e. “system independent”: the composition mechanism
- should not be part of the systems
 - allows to consider **any** system as potentially **open**



They should be

- ▶ CONSERVATIVE
- ▶ FLEXIBLE
- ▶ SAFE

Guaranteeing not to “break” relevant properties of the single systems we compose.

Good composition methods are **safe** when...

If one starts from systems like this....

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If one starts from systems like this....



Good composition methods are **safe**

...one does not end up with something like that

Good composition methods are **safe**

...one does not end up with something like that



The “participants-as-interfaces” (PaI) approach

For systems with message-passing interactions

Introduced (as far as we know) in

FRANCO BARBANERA, UGO DE’LIGUORO, ROLF HENNICKER,
Connecting open systems of communicating finite state machines.
J. Log. Algebraic Methods in Program. (2019)

The “participants-as-interfaces” (PaI) approach

applicable when

participant behaviour



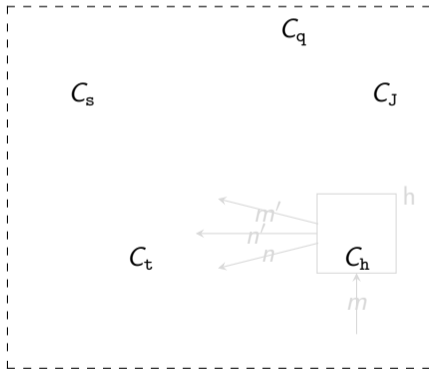
system interface

where

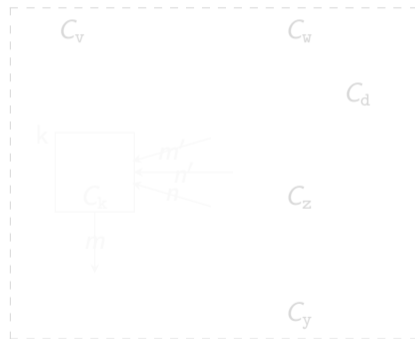
interface = description of the possible interactions with the environment (typically outer systems)

The “participants-as-interfaces” (PaI) approach

S₁

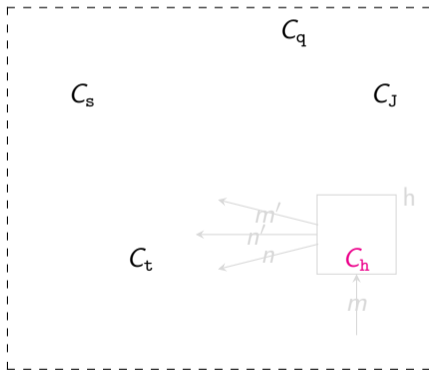


S₂

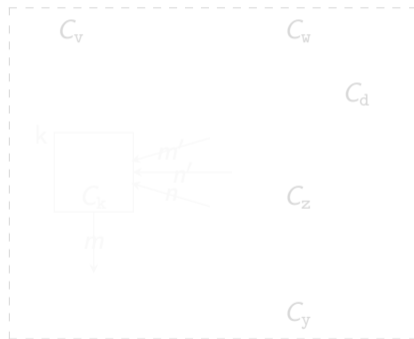


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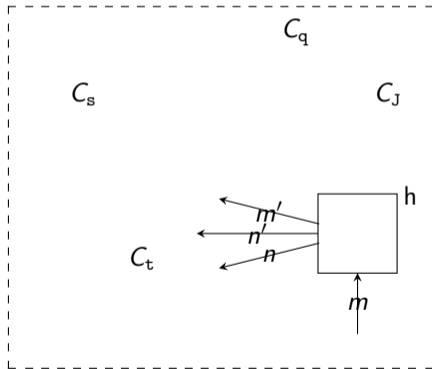


S₂



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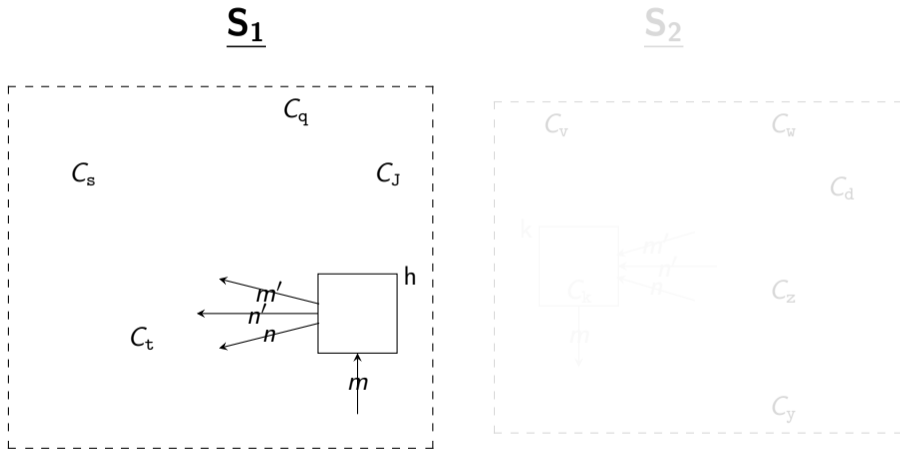
S₁



S₂

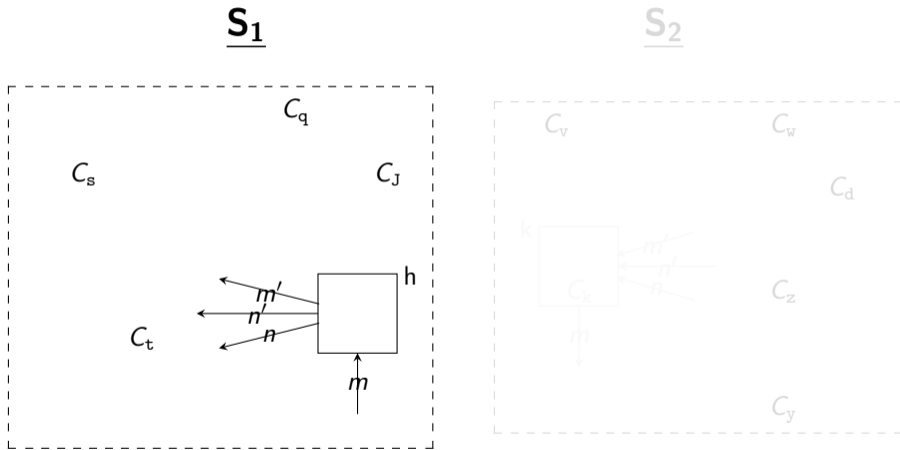


The “participants-as-interfaces” (PaI) approach



We abstract here from the way communications are performed and from the logical order of the exchanged messages.

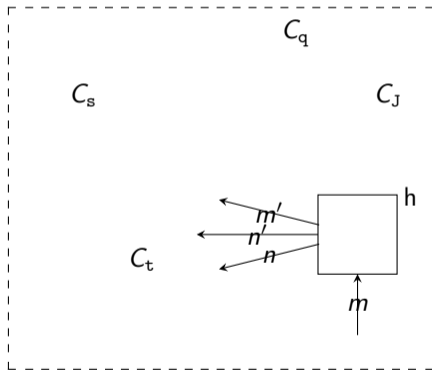
The “participants-as-interfaces” (PaI) approach



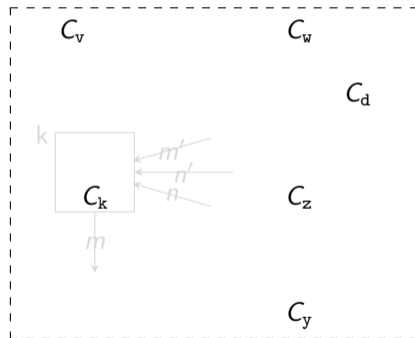
C_h 's behaviour can be looked at as an interface (i.e. a description of what can be offered by an outer system)

The “participants-as-interfaces” (PaI) approach

S₁

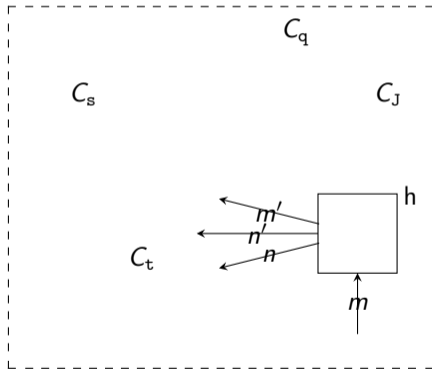


S₂

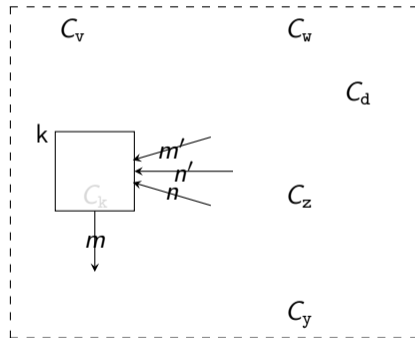


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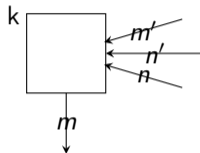
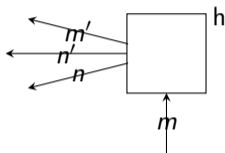
S₁



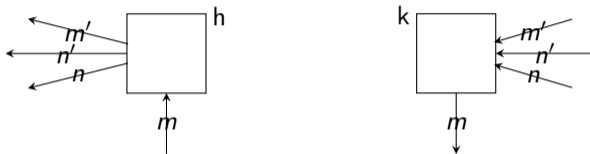
S₂



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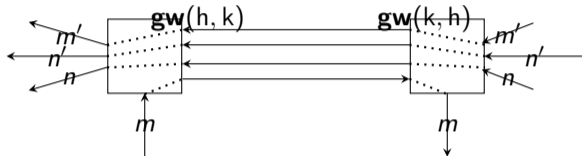


The “participants-as-interfaces” (PaI) approach



COMPATIBLE: an h's input is a k's output, and vice versa

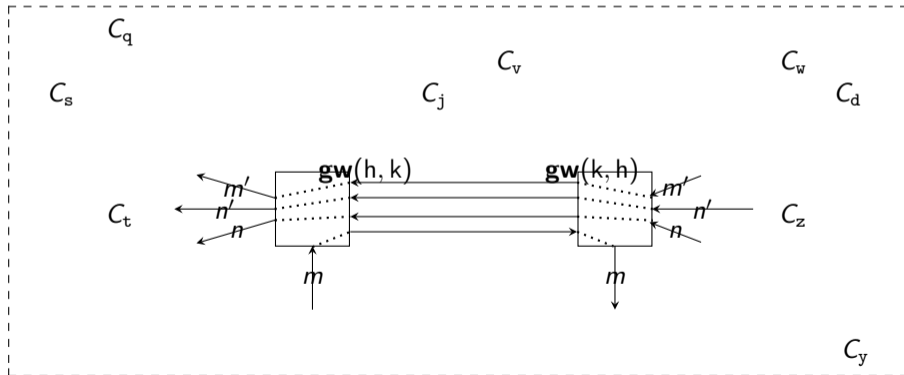
The “participants-as-interfaces” (PaI) approach



Composition via gateways (forwarders)

The “participants-as-interfaces” (PaI) approach

$$\underline{S_1}^{h \leftrightarrow k} \underline{S_2}$$



The “participants-as-interfaces” (PaI) approach

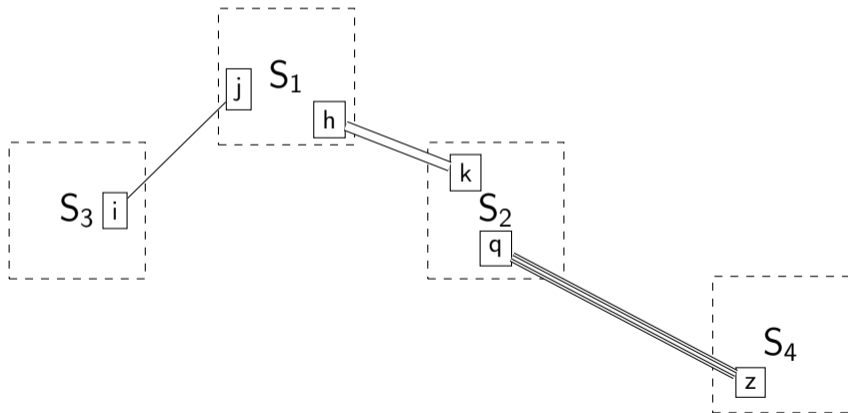
Drawback of binary composition:

By connecting systems two-by-two we get only tree-topologies.

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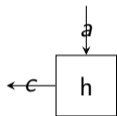
PaI Multicomposition

The composition via gateways trivially extends to simultaneous multiple system composition.

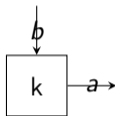
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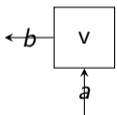
S_1



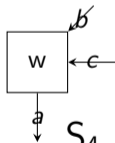
S_2



S_3

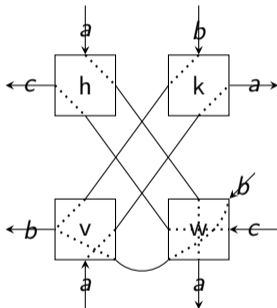


S_4



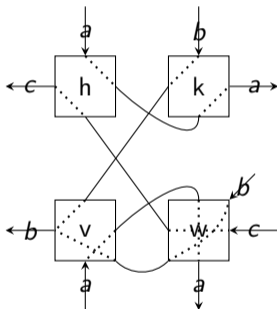
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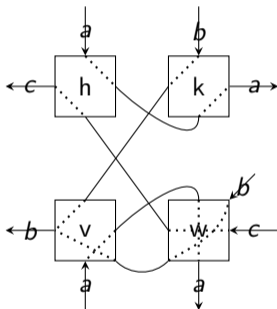
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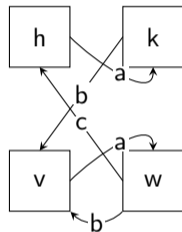
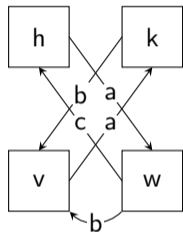
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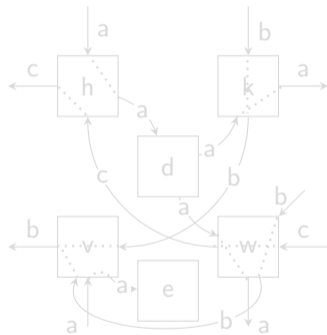
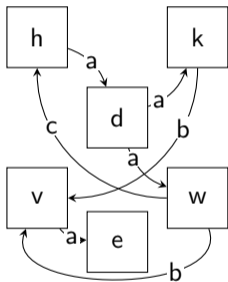
Issues: • Many different “connection policies” (all safe?).

Connection policies as systems

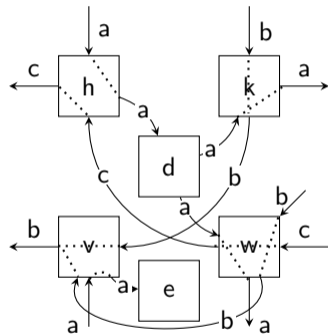
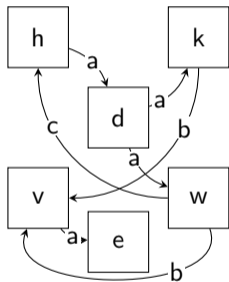


Gateways are built out of the interface participants and a chosen connection policy

Orchestrated connection policies

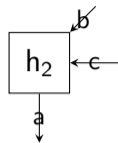
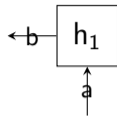


Orchestrated connection policies



Pal is quite general - not relying on any “duality” on interfaces

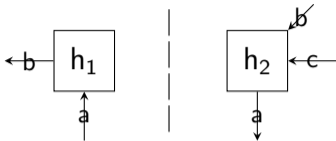
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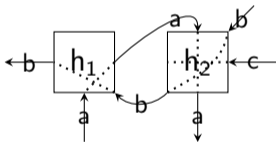
S_2

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S_1

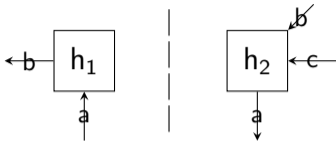


S_2

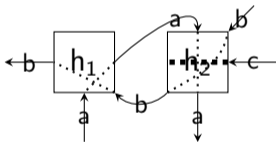


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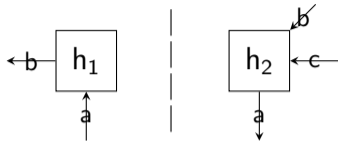


S_2

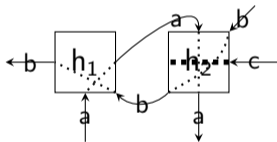


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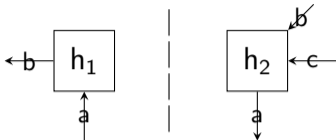
S_2



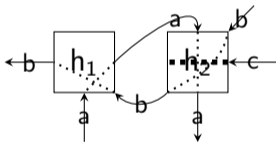
Some properties could be preserved by composition.

Pal is quite general - not relying on any “duality” on interfaces

S_1



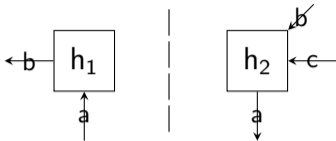
S_2



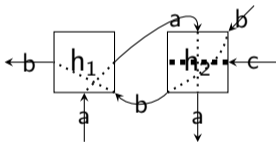
But the interface participants look nonetheless **badly chosen**.

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S_1



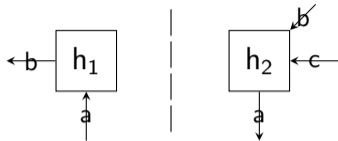
S_2



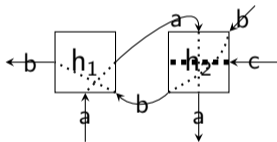
But the interface participants look nonetheless **badly chosen**. Really?

Pal is quite general - not relying on any “duality” on interfaces

S_1



S_2



But the interface participants look nonetheless **badly chosen**. Really?

Perhaps the message c simply does not need to be forwarded.

SISTEMA DOMOTICO

Comfort, efficienza, benessere: tutto sotto controllo.



Illuminazione



Clima



Tapparelle



Sicurezza



Multimedia



DEVICE 2in1

Un unico dispositivo,
due funzioni per il massimo
comfort e benessere.



FUNZIONE 1 CONDIZIONATORE

Regola la temperatura
e raffresca l'ambiente
in modo intelligente
ed efficiente.



FUNZIONE 2 UMIDIFICATORE

Mantiene il livello di umidità
ideale per il benessere
respiratorio e il comfort
in ogni stagione.



SISTEMA DI UMIDIFICAZIONE DISTRIBUITO IN TUTTA LA CASA

Umidità ideale in ogni ambiente, in ogni stagione.



BENESSERE TOTALE

Mantiene il giusto livello di umidità per il comfort e la salute di tutta la famiglia.



PROTEZIONE DELLA CASA

Previene secchezza di legno, crepe nei muri e accumulo di elettricità statica.

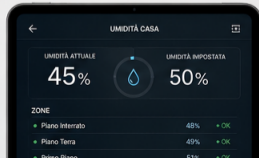


EFFICIENZA E SILENZIOSITÀ

Tecnologia avanzata, consumi ridotti e funzionamento silenzioso.

COME FUNZIONA

- 1 Unità di produzione genera vapore umidificato.
- 2 Il vapore viene immesso nella rete di distribuzione.
- 3 Ugelli di erogazione in ogni ambiente rilasciano il giusto livello di umidità.
- 4 Sensori monitorano l'umidità e regolano automaticamente il sistema.



SOTTOTETTO

Umidità controllata anche negli ambienti più isolati.



SECONDO PIANO

Comfort costante in camere e bagni.



PRIMO PIANO

Ogni stanza riceve la giusta umidità.



PIANO TERRA

Benessere in tutta la zona giorno.



PIANO INTERRATO

Unità di produzione vapore e controllo del sistema.

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Sicurezza

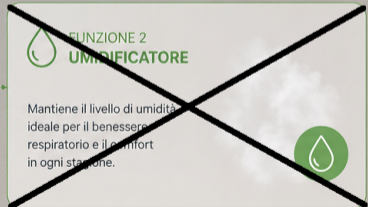


Multimedia



DEVICE 2in1

Un unico dispositivo,
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SISTEMA DOMOTICO

Comfort, efficienza, benessere: tutto sotto controllo.



Illuminazione



Clima



Tapparelle



Sicurezza



Multimedia



partial
GW

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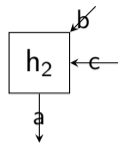
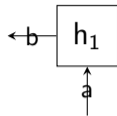
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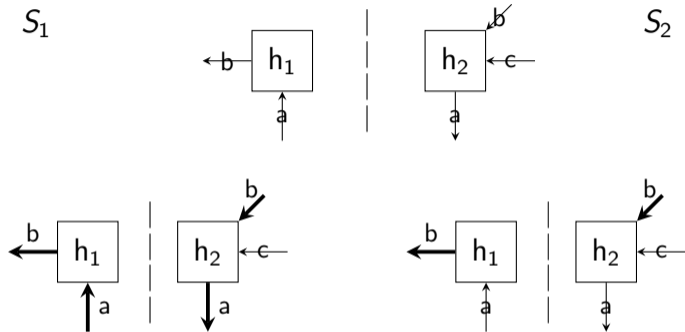
Pal Composition via **partial** gateways - ICE'26

S_1



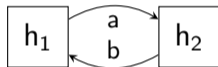
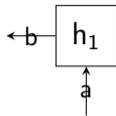
S_2

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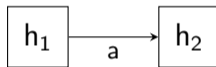
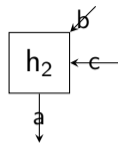


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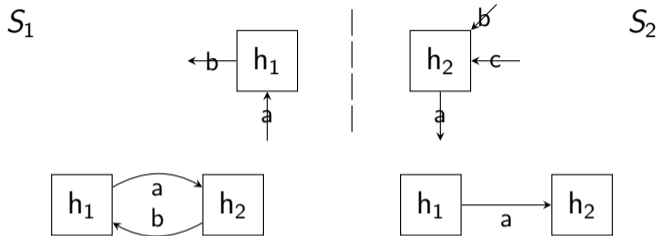
S_1



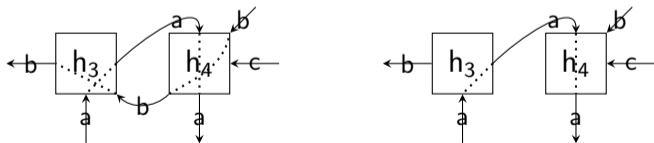
S_2



Pal Composition via **partial** gateways - ICE'26



The above connection policies lead to the following **partial gateways**



The “participants-as-interfaces” (PaI) approach

▶ CONSERVATIVE ✓

▶ FLEXIBLE ✓

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It depends!

▶ SAFE



It depends!

- On the concurrent-system description formalism

▶ SAFE



It depends!

- On the concurrent-system description formalism
- On the communication model

Rather severe restrictions in both

▶ **MultiParty Session Types**

- F. BARBANERA, M. DEZANI-CIANCAGLINI, I. LANESE, E. TUOSTO:
Composition and decomposition of multiparty sessions. JLAMP (2021)
- F. BARBANERA, M. DEZANI-CIANCAGLINI, L. GHERI, N. YOSHIDA:
Multicompatibility for Multiparty-Session Composition. PPDP 2023

▶ **synchronous** version of **CFSM**

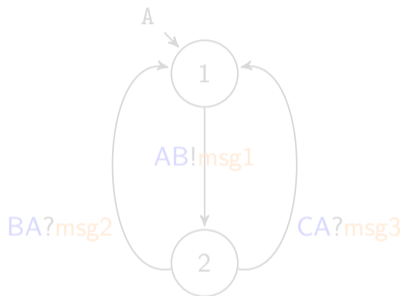
- F. BARBANERA, IVAN LANESE, EMILIO TUOSTO:
Composition of synchronous communicating systems. JLAMP (2023)

Exploiting the PaI approach
in
(**asynchronous**)
Communicating **F**inite **S**tate **M**achines

Communicating Finite State Machines (CFSMs)

A formalism for the description and the analysis of distributed systems.

A machine M_A

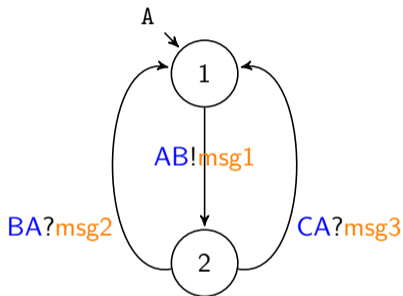


- ▶ M_A can send `msg1` to machine M_B ; **asynchronously**; through the directed buffered FIFO channel `AB`
- ▶ Then, either `msg2` or `msg3` can be received from M_B or M_C ; through channels `BA` or `CA`;
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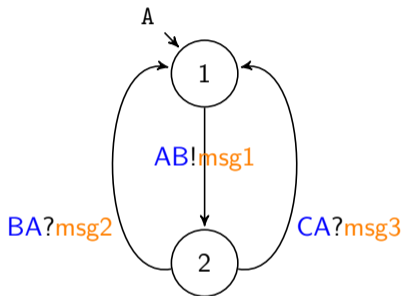


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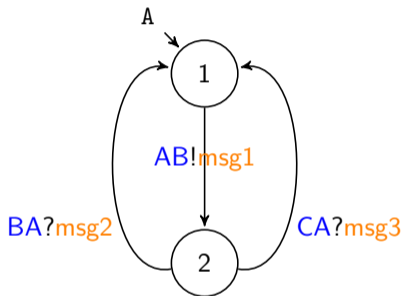


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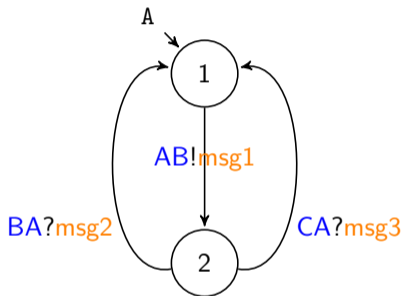


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- Let
- $\{\mathbf{S}_i\}_{i \in I}$ be a set of CFSM systems
 - \mathbb{K} a communication policy for chosen interfaces
 - \mathcal{P} be a communication property

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- ▶ Lock-freedom ✗

Safeness for Pal Orchestrated Multicomposition [Barbanera-Hennicker]

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 - \mathbb{K} an orchestrated comm. policy for chosen interfaces
 - \mathcal{P} be a communication property

IF all S_i 's and \mathbb{K} enjoy \mathcal{P} (and *no-mixed-state* holds)

THEN $\text{OrchMcomp}(\{S_i\}_{i \in I}, \mathbb{K})$ enjoys \mathcal{P} .

WHEN \mathcal{P} is

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- ▶ Orphan-message freedom ✓
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- ▶ Progress ✓
- ▶ Lock-freedom ✗

Safeness for Pal composition via partial gateways - ICE'26

- Let
- $\{S_i\}_{i \in I}$ be a set of CFSM systems
 - \mathbb{K} an orchestrated comm. policy for chosen interfaces and chosen interface transitions.
 - \mathcal{P} be a communication property

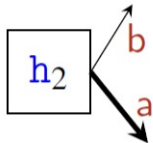
IF all S_i 's and \mathbb{K} enjoy \mathcal{P} (and *no-mixed-state* holds, **even indirect**)

THEN $\text{PartGWcomp}(\{S_i\}_{i \in I}, \mathbb{K})$ enjoys \mathcal{P} .

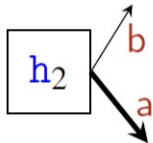
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A requirement needed on interfaces for Pal composition with partial gateways

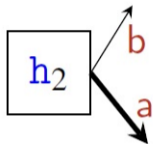


A requirement needed on interfaces for Pal composition with partial gateways



The resulting partial gateways either send **b** or expect **a** to forward it.

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A mixed choice: harmful in CFSM. Property-preservation is disrupted.

All proofs by contradiction

Based on

the notion of **projection**: configurations of original systems and of connection policy obtainable from a configuration of the composed system;

and on the fact that projection preserves (roughly) reachability

All preserved properties are safety properties.

So - roughly - problematic reachable configurations of original systems and of connection policy obtainable from a problematic reachable configuration of the composed system

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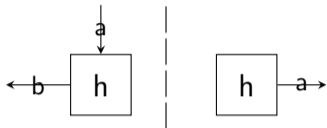
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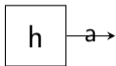
The basis of all:

Partial Fusion Composition - ICE'26

S_1



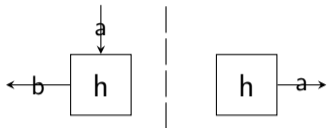
S_2



The basis of all:

Partial Fusion Composition - ICE'26

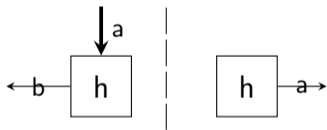
S_1



S_2

If some transitions of h in S_1 perfectly complement the transitions of h in S_2 ...

S_1

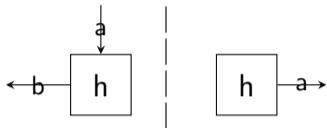


S_2

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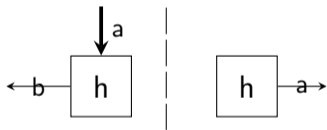
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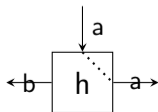
S_1



S_2

... then FUSE the two h 's

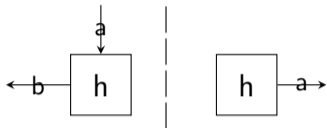
$S_1 \stackrel{h}{\text{fuse}} S_2$



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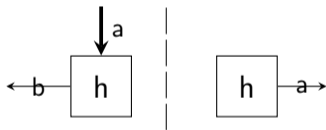
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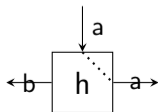
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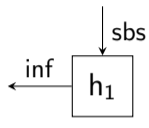
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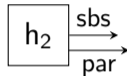
Partial Fusion is **SAFE**

An example of Pal composition via partial gateways

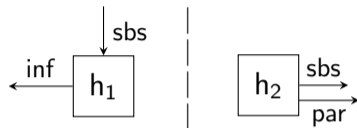
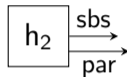
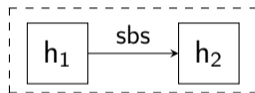
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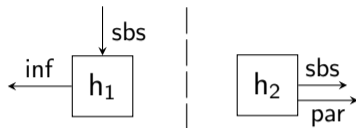
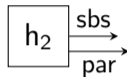
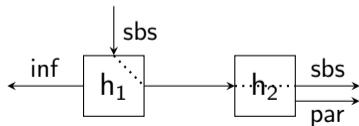
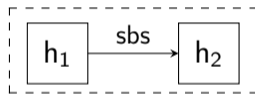
S_2



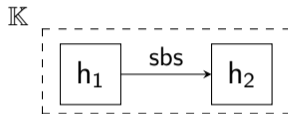
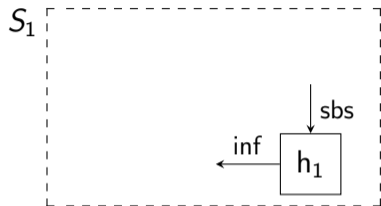
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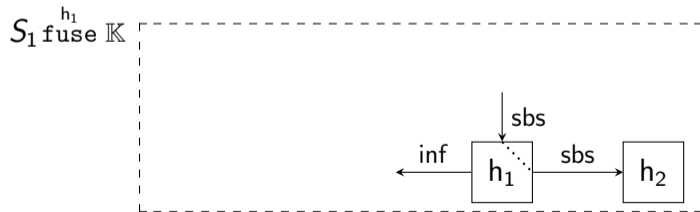
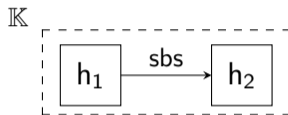
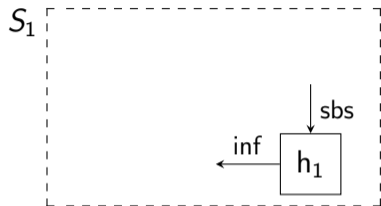
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Binary Pal via partial gateways from Partial Fusion

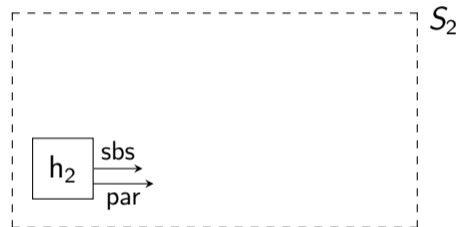
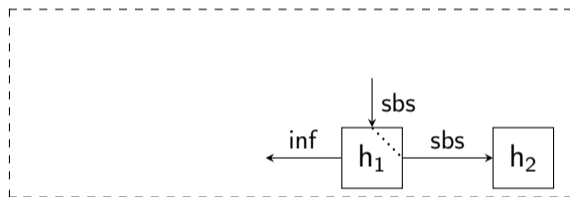


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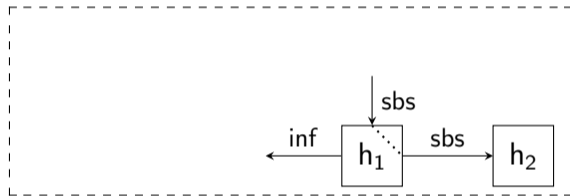
Binary Pal via partial gateways from Partial Fusion

$S_1^{h_1} \text{fuse } \mathbb{K}$

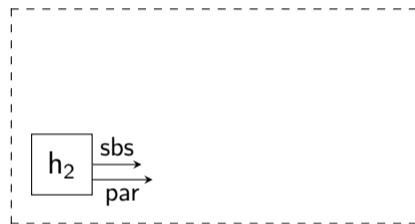


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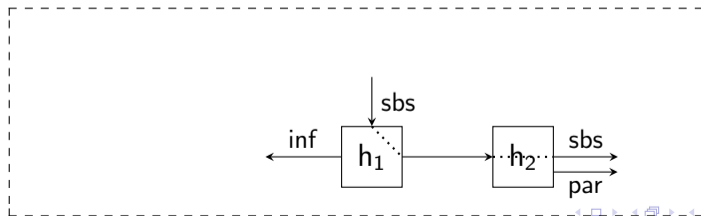
$S_1 \text{ fuse }^{h_1} \mathbb{K}$



S_2



$(S_1 \text{ fuse }^{h_1}) \text{ fuse }^{h_2} S_2$



TAKE-AWAY MESSAGE

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The idea of the paper indirectly stemmed from an observation of a JLAMP reviewer.

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he/she can always be helpful and inspiring

In the future

- ▶ formally show that all the partial Pal multicompositions are obtainable from fusion composition:
- ▶ some (liveness?) conditions to get Lock-freedom preservation for partial fusion composition (and hence lock-freedom for all the partial versions of multicomposition.
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Thanks for your attention

